Main Examination period 2018
MTH6115: Cryptography
Duration: 2 hours

Apart from this page, you are not permitted to read the contents of this question paper until instructed to do so by an invigilator.

## You should attempt ALL questions. Marks available are shown next to the questions.

Only non-programmable calculators that have been approved from the college list of non-programmable calculators are permitted in this examination. Please state on your answer book the name and type of machine used.

Complete all rough work in the answer book and cross through any work that is not to be assessed.

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Exam papers must not be removed from the examination room.

Examiners: W.H.Mannan, B. Noohi

All keys, plaintext and ciphertext will be given in SMALL CAPS, without quotation marks. We use the standard conversion:

| A | B | C | D | E | F | G | H | I | J | K | L | M | N | O | P | Q | R | S | T | U | V | W | X |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 1 | Y | Z |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

When working modulo $n, T_{r}(u)$ will denote $u^{r} \in \mathbb{Z} / n \mathbb{Z}$.
For full marks show all working.

## Question 1. [20 marks]

(a) Encrypt the word can using the affine cipher: $f(x) \equiv 5 x-3 \bmod 26$.
(b) Alice uses the encryption function $g(x) \equiv 4 x-8 \bmod 26$ to communicate with Bob (so he knows that she uses $g$ as her encryption function). Alice requests Bob to bring her a can, encrypting the word CAN with the encryption function $g$.

What might he bring her instead of a can and what feature of the function $g$ would have created this confusion?
(c) The affine cipher $h(x) \equiv 509 x-83 \bmod 1023$ is used to encrypt a message. Find the function which would be used to decrypt the message.
(d) The following ciphertext has been encrypted using an affine cipher. Assume the second most frequent letter in the plaintext is M. Make an educated guess for the most frequent letter in the plaintext. Hence find the decryption function and decrypt the message.

## Question 2. [20 marks]

In all substitution tables the columns correspond to key entries and the rows correspond to plaintext entries.
(a) Encrypt the plaintext cat with the key TCG using the substitution table $S$ :

$$
S=\quad \begin{array}{c|cccc} 
& \mathrm{A} & \mathrm{C} & \mathrm{G} & \mathrm{~T} \\
\hline \mathrm{~A} & \mathrm{~T} & \mathrm{G} & \mathrm{C} & \mathrm{~A} \\
\mathrm{C} & \mathrm{~A} & \mathrm{C} & \mathrm{~A} & \mathrm{~T} \\
\mathrm{G} & \mathrm{G} & \mathrm{~A} & \mathrm{C} & \mathrm{G} \\
\mathrm{~T} & \mathrm{C} & \mathrm{~T} & \mathrm{G} & \mathrm{C}
\end{array}
$$

(b) Give an example of a plaintext message, which when encrypted with key TCG and substitution table $S$, cannot be decrypted even by the intended recipient (who is in possession of both the key and substitution table).
(c) The substitution table $T$ is given by:

$$
T=\quad \begin{array}{c|cccc} 
& \text { A } & \text { C } & \mathrm{G} & \mathrm{~T} \\
\hline \mathrm{~A} & \mathrm{~T} & \mathrm{~T} & \mathrm{~T} & \mathrm{G} \\
\mathrm{C} & \mathrm{G} & \mathrm{C} & \mathrm{~A} & \mathrm{~T} \\
\mathrm{G} & \mathrm{~A} & \mathrm{~A} & \mathrm{C} & \mathrm{~A} \\
\mathrm{~T} & \mathrm{C} & \mathrm{G} & \mathrm{G} & \mathrm{C}
\end{array}
$$

What feature of $T$ means the problem in (b) will not occur?
(d) A plaintext message $M$ is encrypted using substitution table $T$. Initially you know that:

$$
\mathbb{P}(M=\mathrm{TAT})=0.1, \quad \mathbb{P}(M=\mathrm{GAG})=0.3, \quad \mathbb{P}(M=\mathrm{ACT})=0.6
$$

You intercept the ciphertext ctc. Modify the above probabilities accordingly.
(e) In addition to your answer to (c), what further feature must a substitution table have in order to be a Latin square? What is the advantage of having this feature? Name the theorem that states that Latin squares have this advantage.

## Question 3. [22 marks]

(a) A shift register is described by the polynomial $x^{3}+x^{2}+1$. Starting with initial configuration 001 calculate the key generated by this register. Is the shift register primitive?
(b) Calculate the number of runs of 0's and the number of runs of 1's of each length in this key. Also calculate the out of phase autocorrelations of each shift and the in phase autocorrelation. With reference to these numbers, determine which of Golomb's postulates this key satisfies.
(c) State a theorem which would have told you immediately which of Golomb's postulates this key satisfies, without performing the calculations in (b).
(d) In a key generated by a primitive $n$-bit shift register, how many runs of 1 's of length $n-1$ are there? Explain your answer without using any results that you do not prove.

Question 4. [16 marks] For full marks any methods you use should have fast (polynomial time) implementations on a classical computer.
(a) Bob is using RSA cryptography with public key ( $n=2035153$, $e=11$ ). Eve hides a camera in Bob's home and catches him typing in his private key $d=1283$. She then intercepts the encrypted number $T_{e}(x) \equiv 5 \bmod 2035153$. Decrypt this to find $x$.
(b) Bob finds the camera and realises that Eve has seen him type in his private key He then switches his public key to $e^{\prime}=13$. Show that this is a mistake by computing his new private key (do not worry about the size of your answer, as long as it works).

[^0]
## Question 5. [22 marks]

(a) A publicly known prime number $p$ is fixed. Alice has a secret number $a$. Write down a linear equation in $x$ which guarantees that $T_{x}$ is the inverse of $T_{a}$. Name the method needed to find $x$ by solving this equation.
(b) Bob also has a secret number $b$ and $T_{y}$ is the inverse of $T_{b}$.

Alice wishes to secretly communicate a number $m \bmod p$ to Bob. Let $m_{0}, m_{4}=m$, and let $m_{1}, m_{2}, m_{3}$ be numbers they exchange publicly in that order. So first $m_{1}$ is sent publicly, then $m_{2}$, then $m_{3}$.

For $i=1,2,3,4$ write down:

- Which of the functions $T_{a}, T_{b}, T_{x}, T_{y}$ is applied to $m_{i-1}$ to get $m_{i}$,
- Who is applying the function,
- Which direction (if any) $m_{i}$ is sent ('Alice to Bob' or 'Bob to Alice').
(c) Eve sees all the numbers being sent publicly. Suppose now it is publicly known that $m$ is a primitive generator modulo $p$. If Eve invented a fast algorithm for solving the discrete logarithm problem, how could she break the encryption and find $m$ ?


## End of Paper.


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