

# WEEK 12 Lecture 1

## CHAPTER 8 THE CONFIGURATION MODEL

It is a model to generate random networks with a given degree sequence  $\leftarrow$  degree distribution

### 8.1 THE CONFIGURATION MODEL

DEF

Considers the ensemble of networks with  $N$  nodes and

with a given degree sequence  $\{K_i\} = \{K_1, K_2, \dots, K_N\}$

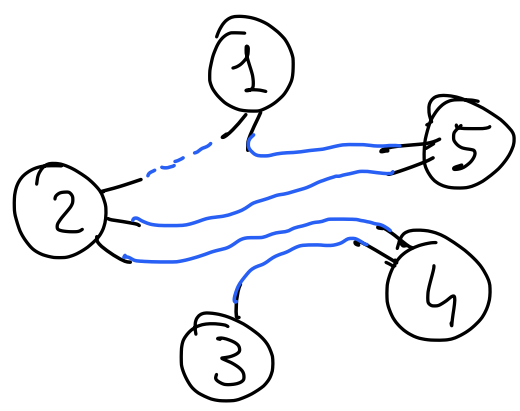
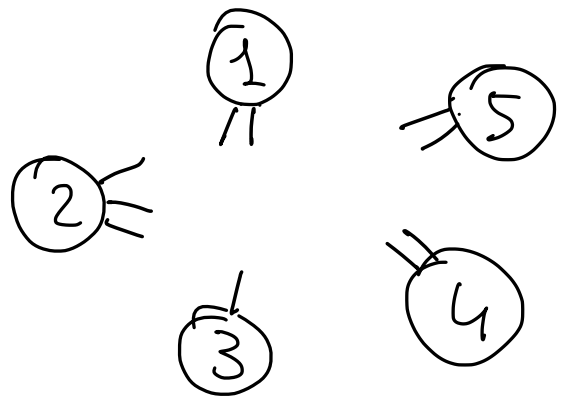
constructed as follows:

- (a) Place  $K_i$  stubs on node  $i$

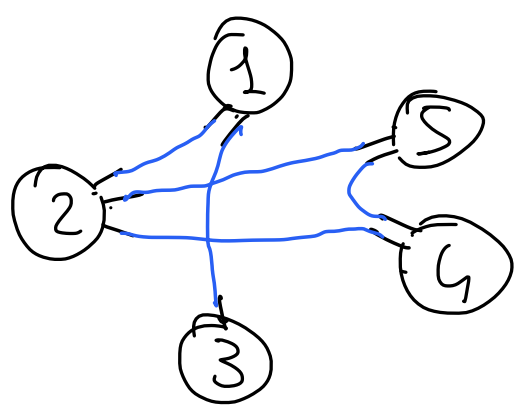
$$\sum_i K_i = 2L$$

- Ⓛ Match each stub randomly with another stub
- ⓐ When each stub is matched, if the network does not contain tadpoles or multiple links then stop, otherwise start again with Ⓛ

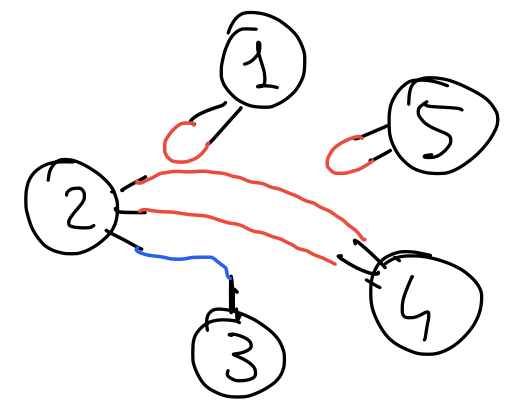
Degree sequence  $\left\{ \begin{matrix} k_1 & k_2 & k_3 & k_4 & k_5 \\ 2, & 3, & 1, & 2, & 2 \end{matrix} \right\}$   $N=5$  nodes



ACCEPT



ACCEPT



REJECT

# 8.2 UNCORRELATED NETWORKS

## DEF

A network with degree sequence  $\{k_i\}$  is UNCORRELATED if

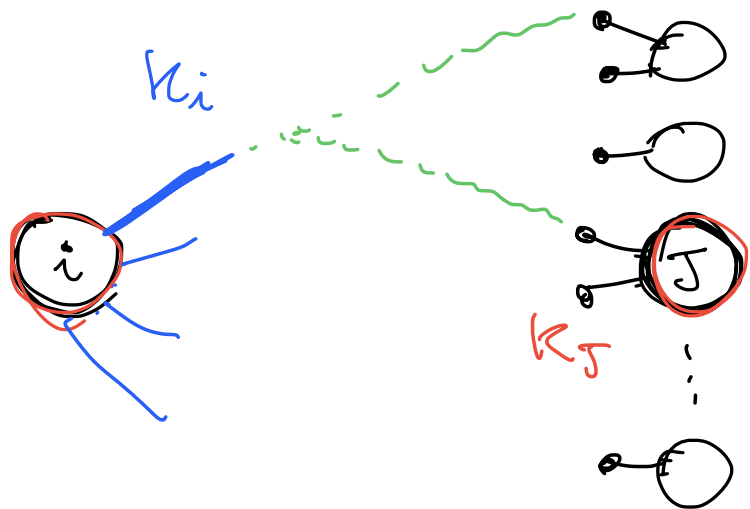
Probability that node  $i$  is connected to node  $j$

$$P_{ij} = \frac{k_i k_j}{\langle k \rangle N}$$

Probability that a link from node  $i$  connects node  $j$

$$q_{ij} = \frac{k_j}{\langle k \rangle N} \equiv q_j$$

uncorrelated if  $q_{ij}$  only depends on  $j$



$k_j$  possibilities to connect a link of  $i$  to node  $j$ , out of  $\sum_j k_j = 2L$

$$\langle k \rangle = \frac{2L}{N}$$

$$= \langle k \rangle N$$

(3)

Hence  $q_{ij} = q_j = \frac{k_j}{\langle k \rangle N}$  and  $p_{ij} = k_i \cdot q_{ij} = \frac{k_i k_j}{\langle k \rangle N}$

## DEF | STRUCTURAL CUTOFF

The structural cutoff of a network is the maximum allowed degree given by

$$K_s = \sqrt{\langle k \rangle N}$$

A network has a structural cutoff if  $k_i \leq K_s \quad \forall i$

## PROPOSITION

The networks generated by the configuration model are **UNCORRELATED** iff they have a **STRUCTURAL CUTOFF**

Proof (of the necessary condition only)

a network is uncorrelated  $\Rightarrow$  the network has the structural cutoff

$$\downarrow$$
$$p_{ij} = \frac{k_i k_j}{\langle k \rangle N}$$

Let us indicate as  $\bar{K}$  the largest degree in the network and assume

$$k_i = k_j = \bar{K}$$

prob that  $i$  and  $j$  are connected  $\rightarrow$   $P_{ij} = \frac{\bar{K} \cdot \bar{K}}{\langle k \rangle N}$

by definition

$$P_{ij} \leq 1$$

$$\frac{\bar{K}^2}{\langle k \rangle N} \leq 1$$

$$\bar{K} \leq \sqrt{\langle k \rangle N}$$

Since  $\bar{K} = \max_i k_i$  we then have  $k_i \leq \bar{K}_s \quad \forall i$

## PROPOSITION

In an uncorrelated network the probability  $q(k)$  that following a link we reach a node of degree  $k$  is

$$q(k) = \frac{k P(k)}{\langle k \rangle}$$

degree distribution

Proof

$$q_{ij} = q_j = \frac{k_j}{\langle k \rangle N}$$

Hence  $\frac{k}{\langle k \rangle N}$  is the probability that following a link we get into a node  $j$  having  $k_j = k$

probability that following a link we get in nodes of degree  $k$

$$q(k) = \frac{k}{\langle k \rangle N} \cdot (\# \text{ of nodes of degree } k)$$

$$P(k) = \frac{N(k)}{N}$$

$$N(k) = P(k) \cdot N$$

$$q(k) = \frac{k}{\langle k \rangle N} \cdot P(k) \cdot N = \frac{k P(k)}{\langle k \rangle}$$

The average node degree is  $\langle k \rangle$ , while the average degree of the neighbours of a node is NOT  $\langle k \rangle$

### PROPOSITION

The average degree  $k_{nn}$  of the neighbours of a node is given by

$$k_{nn} = \frac{\langle k^2 \rangle}{\langle k \rangle}$$

nearest-neighbours

Proof

$$\langle k \rangle = \sum_k k p(k)$$

$$\langle k^2 \rangle = \sum_k k^2 p(k)$$

$$k_{nn} = \sum_k k q(k) = \sum_k k \frac{k p(k)}{\langle k \rangle} = \frac{1}{\langle k \rangle} \sum_k k^2 p(k) = \frac{\langle k^2 \rangle}{\langle k \rangle}$$

$$q(k) = \frac{k p(k)}{\langle k \rangle}$$

See example in FORM. ASSIGN 10

**FRIENDSHIP PARADOX**: your friends have more friends than you have!!

$\langle k \rangle$  = av # of friends of an individual

$k_{nn} = \frac{\langle k^2 \rangle}{\langle k \rangle}$  = av. # of friends of the neighbours of an individual

We have  $k_{nn} > \langle k \rangle$  when  $\frac{\langle k^2 \rangle}{\langle k \rangle} > \langle k \rangle$

$$\langle k^2 \rangle > \langle k \rangle^2$$

which is always true with the exception

of regular networks where  $k_i = \langle k \rangle \forall i$

$$\langle (k - \langle k \rangle)^2 \rangle > 0$$

# 8.4 CLUSTERING COEFFICIENT

## PROPOSITION

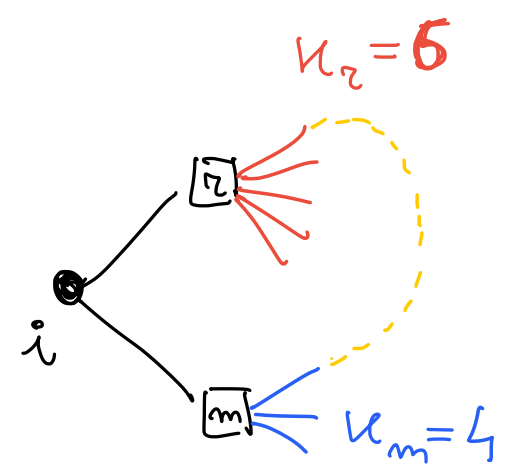
The clustering coefficient of an uncorrelated network from the configuration model is:

$$C = \frac{\langle k \rangle}{N} \left[ \frac{\langle k^2 \rangle - \langle k \rangle^2}{\langle k \rangle^2} \right]^2$$

Proof

$$C_i = \frac{T_i}{\frac{k_i(k_i-1)}{2}} = \text{fraction of pairs of neighbours of node } i \text{ that are linked}$$

probability that 2 neighbours of a node are linked



$n$  with degree  $k_n \rightarrow k_n - 1$  remaining links

$m$  with degree  $k_m \rightarrow k_m - 1$  remaining links

Since in the configuration model stubs are matched at random, then the probability that  $n$  is linked to  $m$  is

$$P_{nm} = \frac{(k_n - 1)(k_m - 1)}{\langle k \rangle N}$$



$r$  has degree  $k_r$  with a probability  $q(k_r)$   
 $m$  "  $k_m$  " "  $q(k_m)$

Hence averaging over all possible values of  $k_r$  and  $k_m$

$$C = \sum_{k_r} \sum_{k_m} q(k_r) q(k_m) \cdot \frac{(k_r-1)(k_m-1)}{\langle k \rangle N} =$$

$$= \frac{1}{\langle k \rangle N} \left( \sum_{k_r} q(k_r) (k_r-1) \right) \cdot \left( \sum_{k_m} q(k_m) (k_m-1) \right)$$

$$= \frac{1}{\langle k \rangle N} \left[ \frac{\langle k(k-1) \rangle}{\langle k \rangle} \right]^2 =$$

$$= \frac{\langle k \rangle}{N} \left[ \frac{\langle k^2 \rangle - \langle k \rangle}{\langle k \rangle^2} \right]^2$$

Poiss

A connection term that can become large if  $P(k)$  is scale-free which can lead to a finite  $C$

notice!

$$\sum_m \frac{k_m P(k_m)}{\langle k \rangle} (k_m-1) =$$

$$= \frac{1}{\langle k \rangle} \sum_m k_m (k_m-1) P(k_m)$$

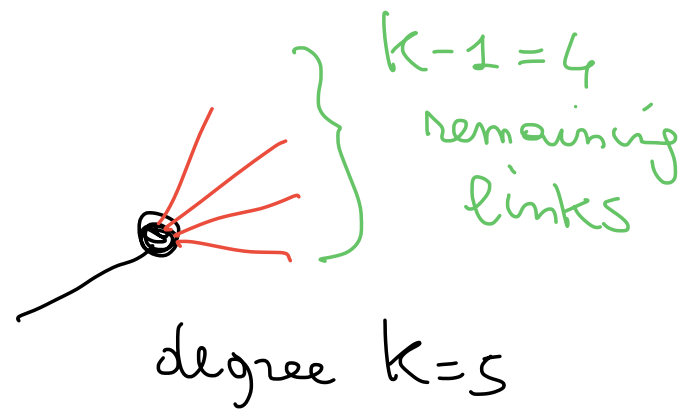
$$= \frac{\langle k(k-1) \rangle}{\langle k \rangle}$$

# 8.5 AVERAGE DISTANCE

## DEF | BRANCHING RATIO

The branching ratio  $b_k$  of a node of degree  $k$  is

$$b_k = k - 1$$



## PROPOSITION

the average branching ratio  $\bar{b}$  of a node reached by following a link in an uncorrelated network is:

$$\bar{b} = \frac{\langle k(k-1) \rangle}{\langle k \rangle}$$

used in Sect 7.6

Proof

$$\bar{b} = \sum_k b_k q(k) = \sum_k (k-1) k \frac{P(k)}{\langle k \rangle} = \frac{1}{\langle k \rangle} \sum_k (k^2 - k) P(k)$$

$q(k) = \frac{k P(k)}{\langle k \rangle}$

$$= \frac{1}{\langle k \rangle} \langle k(k-1) \rangle = \frac{\langle k^2 \rangle - \langle k \rangle}{\langle k \rangle}$$

## PROPOSITION

The average distance  $l$  in an uncorrelated network with finite average branching ratio  $\bar{b}$  is, in the limit  $N \rightarrow \infty$

$$l \approx \frac{\ln N}{\ln \frac{\langle k(k-1) \rangle}{\langle k \rangle}} + \text{Constant}$$

displays the  
S.W.D.P.

Proof

Assuming the network to be "locally tree-like" ( $C \rightarrow 0$  as  $N \rightarrow \infty$ )

We compute the # of nodes at distance  $d$  from a given (generic) node

$$N_d \approx \langle k \rangle \bar{b}^{d-1} = \langle k \rangle \left[ \frac{\langle k(k-1) \rangle}{\langle k \rangle} \right]^{d-1}$$

$\bar{b} = \frac{\langle k(k-1) \rangle}{\langle k \rangle}$   
 average branching ratio

We can now repeat the same argument as in Sect 7.6

where we had

$$N_d \approx \langle k \rangle e^{d-1} \quad \text{and we found} \quad l \approx \frac{\ln N}{\ln e}$$

and we get

$$l \approx \frac{\ln \frac{N}{\langle k \rangle}}{\ln \frac{\langle k(k-1) \rangle}{\langle k \rangle}} + 1 = \frac{\ln N}{\ln \frac{\langle k(k-1) \rangle}{\langle k \rangle}} + \text{Const.}$$

$$N \approx \langle k \rangle \bar{b}^{l-1}$$

$$\ln \frac{N}{\langle k \rangle} = (l-1) \ln \bar{b}$$

$$l = \frac{\ln \frac{N}{\langle k \rangle}}{\ln \bar{b}} + 1$$