

# WEEK 11 Lecture 1

## CHAPTER 4

### SMALL-WORLD NETWORKS

Many real-world networks have at the same time

in W3

A small distances between their nodes (S.W.D.P.)

B large number of triangles

- How to measure this?



- It is easy to induce A+B  
in networks with "many" links

PUZZLE: real networks are SPARSE

shortest-path distance matrix

$$\text{SPM} = \{d_{ij}\}$$

$$\text{DIAMETER } D = \max_{i,j} \{d_{ij}\}$$

CHARACTERISTIC

$$\text{PATH LENGTH } \ell = \langle d_{ij} \rangle$$

Introduction of WS SMALL-WORLD network model

## M.2 CLUSTERING COEFFICIENT

Watts and Strogatz (WS)  
in Nature 1998 paper

### DEF

The NODE CLUSTERING COEFFICIENT  $C_i$  of node  $i$  is

$$C_i = \begin{cases} \frac{T_i}{\frac{k_i(k_i-1)}{2}} & \text{if } k_i > 1 \\ 0 & \text{if } k_i = 0, 1 \end{cases}$$

where  $k_i$  is the degree of node  $i$ , and

$T_i$  is the # of triangles passing through node  $i$

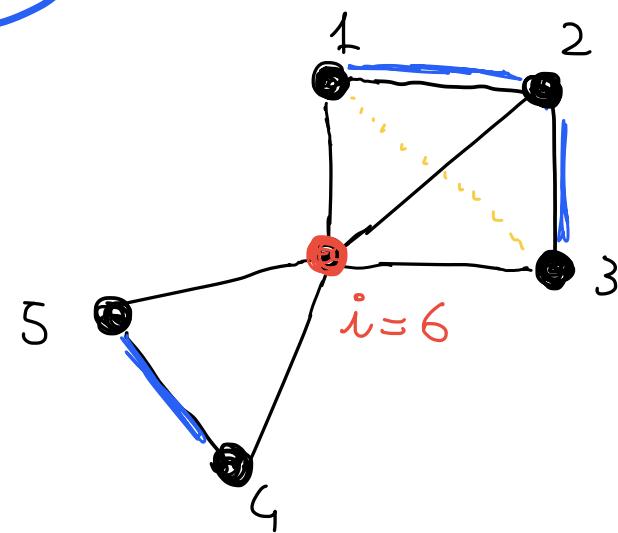
The NETWORK CLUSTERING COEFFICIENT  $C_{ws}$

is  $C_{ws} = \frac{1}{N} \sum_{i=1}^N C_i$  average over all the nodes  
in the network

Notice that  $\frac{k_i(k_i-1)}{2} = \binom{k_i}{2}$  is the max # of triangles  
that can pass through a node of degree  $k_i$

Hence  $C_i \in [0, 1] \quad \forall i \Rightarrow C \in [0, 1]$

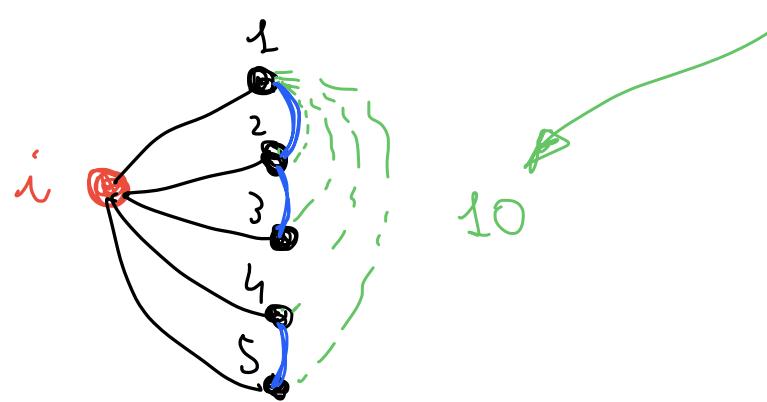
**Ex**



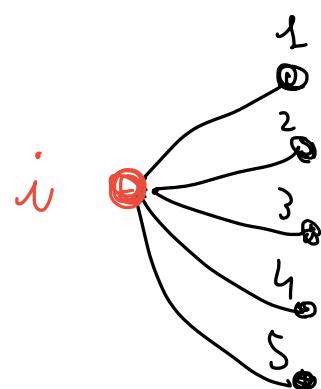
$$N = 6$$

$$k_i = 5$$

$$\binom{k_i}{2} = \frac{k_i(k_{i-1})}{2} = \frac{5 \cdot 4}{2} = 10$$



max # of triangles  
that can pass  
through node  $i=6$



Finally for  $i=6$

$$C_i = \frac{T_i}{\binom{k_i}{2}} = \frac{3}{10}$$

(3)

$$C_1 = \frac{1}{\binom{1}{2}}$$

$$\binom{2}{2} = 1$$

$$C_2 = \frac{2}{3}$$

$$\binom{3 \cdot 2}{2}$$

$$C_3 = \frac{1}{1}$$

$$C_4 = \frac{1}{1}$$

$$C_5 = \frac{1}{1}$$

$$C_6 = \frac{3}{10}$$

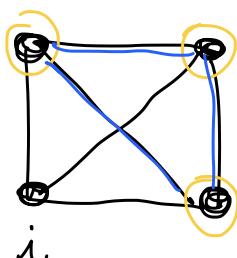
$$C_{ws} = \frac{1}{N} \sum_{i=1}^N C_i = \frac{1}{6} \left[ 1 + \frac{2}{3} + 1 + 1 + 1 + \frac{3}{10} \right] = \frac{1}{6} \cdot \frac{149}{30} \approx 0.83..$$

**Ex**

### COMPLETE NETWORKS

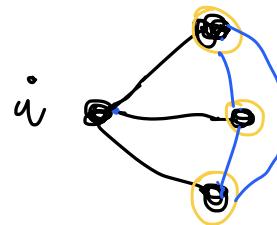
All the nodes have clustering coefficient equal to 1 :  $C_i = 1 \forall i$

$IK_4$



$$N = 4 \\ k_i = 3 \quad \forall i$$

$$C_i = \frac{3}{\frac{3 \cdot 2}{2}} = 1 \quad \forall i$$



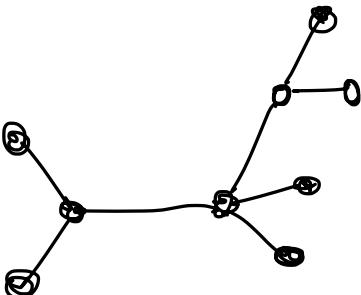
Complete netw  
 $C_{ws}$

$$= 1$$

(3)

Ex

TREES



$$C_i = 0 \quad \forall i \Rightarrow C_{ws}^{\text{tree}} = 0$$

A because a tree has no cycles of any size (so no triangles)



linear chain

There are alternative measures to  $C_{ws}$

DEF

a.k.a. TRANSITIVITY

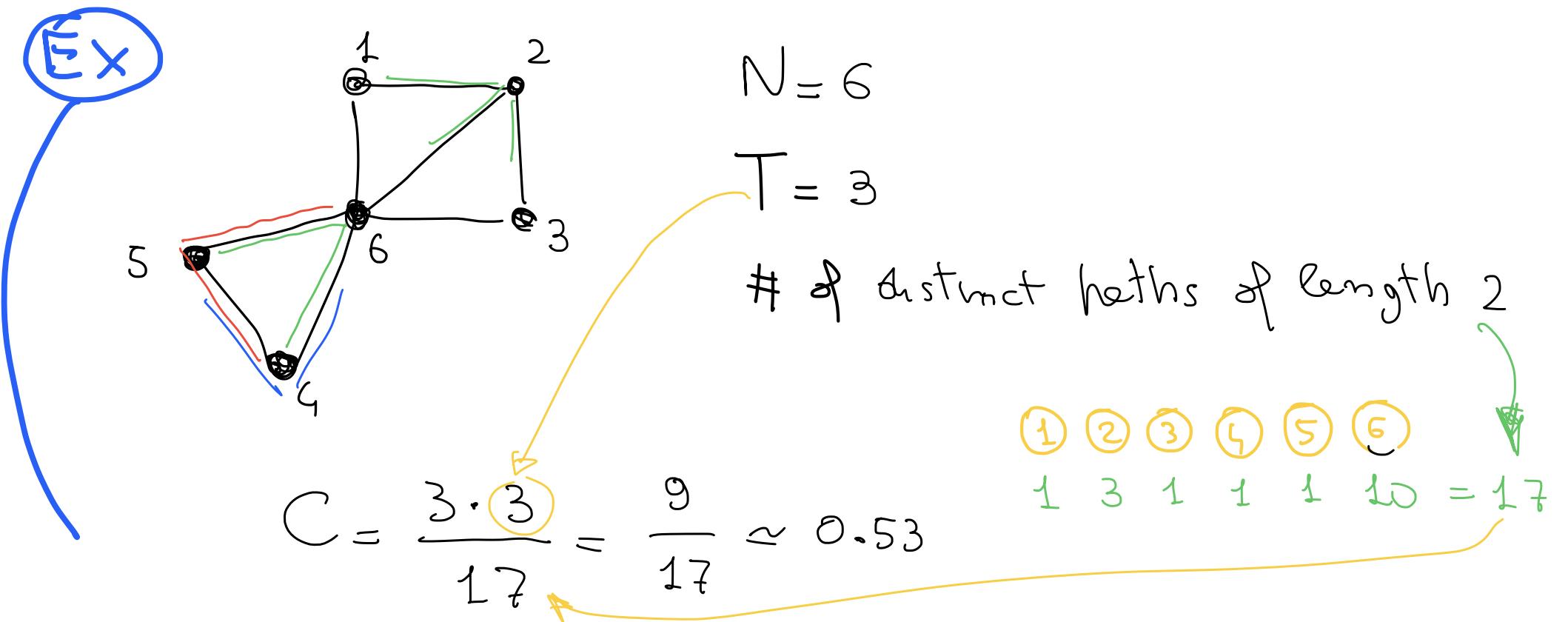
introduced in the social network community

The GLOBAL CLUSTERING COEFFICIENT

C

$$C = \frac{3 \cdot T}{\# \text{ of distinct paths of length 2}}$$

where T is the # of triangles in the network



### 7.3 SMALL-WORLD NETWORKS

**DEF** SMALL-WORLD (DISTANCE) PROPERTY (S.W.D.P.)

A network has the SWDP if

$$D \approx O(\ln N)$$

big O

or if

the diameter is of the same order of magnitude as  $\ln N$

$D \simeq O(\ln N)$  is of a smaller order of magnitude

small O

$$D \simeq O(\ln N) \Rightarrow \lim_{N \rightarrow \infty} \frac{D}{\ln N} = c < \infty \quad \text{with } c > 0$$

$$D \simeq \underline{O}(\ln N) \Rightarrow \lim_{N \rightarrow \infty} \frac{D}{\ln N} = 0$$

Basically SWDP  $\Leftrightarrow$

$$\lim_{N \rightarrow \infty} \frac{D}{\ln N} = \text{constant} < \infty$$

## PROPOSITION

If a network has the SWDP. then either

$$l \simeq O(\ln N) \quad \text{or} \quad l \simeq o(\ln N)$$

characteristic heth length  
(average distance)

Hence

$$\lim_{N \rightarrow \infty} \frac{l}{\ln N} < \infty$$

Proof:  $l \leq D$

## REMARK 1

We will later prove that

$$l^{\text{Rand}} = \frac{\ln N}{\ln \langle k \rangle}$$

Hence SWDP means that  $l$  of a network is of the same order of magnitude of the average distance  $l^{\text{rand}}$  in a random network having the same average degree  $\langle k \rangle$ .

## DEF

A network has a HIGH CLUSTERING COEFFICIENT  $C_{ws}$  if

$$C_{ws} \gg$$

$$\frac{\langle k \rangle}{N}$$

## REMARK 2

We will later prove that

$$C_{ws}^{\text{Rand}} = \frac{\langle k \rangle}{N}$$

## DEF

### SMALL-WORLD NETWORKS

Networks with

[A] the S.W.D.P.

[B] high clustering coefficient

7.5

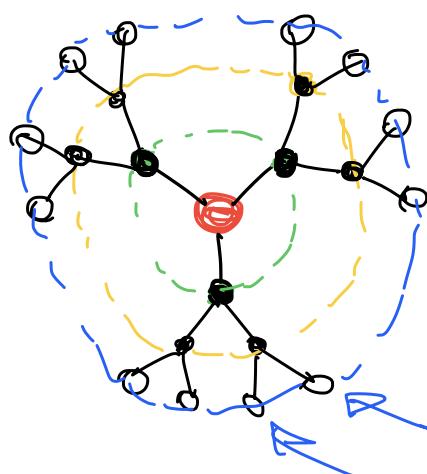
## CAYLEY TREES

### DEF CAYLEY TREE

A symmetric regular tree constructed starting from a given node (the origin) of degree  $K$ . Every node at distance  $0 < d < P$  from the origin has degree  $K$ , while the nodes at distance  $d = P$  have degree 1. The quantity  $b = K - 1$  is called the **BRANCHING RATIO**.

Ex

Cayley tree with  $K = 3$   $P = 3$



branching ratio  
 $b = K - 1 = 2$

leaves (degree = 1)

## PROPOSITION

Cayley trees have the SWDP as:

$$D \approx 2 \frac{\ln N}{\ln b} \quad \text{for } N \gg 1$$

Proof

(a) We first prove that:

# of nodes at distance  $d$  from the origin

$$N_d = \begin{cases} 1 & d=0 \\ k b^{d-1} & 1 \leq d \leq P \end{cases}$$

For  $d=0$

$$N_0 = 1 \quad (\text{this is the origin})$$

For  $d > 0$  Proof of  $N_d = k b^{d-1}$  by induction

- $\boxed{d=1} \rightarrow N_1 = k b^{1-1} = k$  which is true because the origin has  $k$  neighbours

- Then we assume  $N_d = k b^{d-1}$  holds true for  $1 \leq d \leq P$  and we move  $N_{d+1} = k b^d$

(10)

$$N_{d+1} = N_d \cdot b = k b^{d-1} \cdot b = k b^d$$

↑  
each node at distance d  
leads to  $b = k-1$  nodes

(b) We then calculate  $N$  (the total # of nodes in the network) as

$$N = \sum_{d=0}^P N_d = N_0 + \sum_{d=1}^P N_d = 1 + k \sum_{d=1}^P b^{d-1} = 1 + k \sum_{m=0}^{P-1} b^m$$

# of nodes  
at distance  
 $d \leq P$  from the origin

$N_d = k b^{d-1} \quad 1 \leq d \leq P$

Using the GEOMETRIC SUM FORMULA

$$\sum_{n=0}^M r^n = \frac{1 - r^{M+1}}{1 - r}$$

With  $M = P-1$  and  $r = b$  we get:

$$N = 1 + k \frac{1 - b^{P-1+1}}{1 - b} = 1 + k \frac{1 - b^P}{1 - b} = 1 + k \frac{(k-1)^P - 1}{k-2}$$

$\nearrow$   
 $b = k-1$

(c)

We can now calculate the diameter  $D$  as

$$D = 2P$$

(11)

$$N = 1 + k \frac{\left(\frac{k-1}{2}\right)^{\frac{D}{2}} - 1}{k-2} \quad (N-1) \frac{k-2}{k} = \left(\frac{k-1}{2}\right)^{\frac{D}{2}} - 1$$

$$\ln \left[ (N-1) \frac{k-2}{k} + 1 \right] = \frac{D}{2} \ln (k-1)$$

$$D = 2 \frac{\ln \left[ (N-1) \frac{k-2}{k} + 1 \right]}{\ln (k-1)}$$

For  $N \gg 2$

$$\begin{aligned} \ln \left[ (N-1) \frac{k-2}{k} + 1 \right] &\approx \ln \left[ (N-1) \frac{k-2}{k} \right] \approx \ln \left[ N \frac{k-2}{k} \right] = \\ &= \ln N - \ln k + \ln(k-2) \approx \ln N \end{aligned}$$

$$D \approx 2 \frac{\ln N}{\ln(k-1)} = 2 \frac{\ln N}{\ln b}$$

### PROPOSITION

In a Cayley tree  $C_i = 0 \forall i$ . Hence  $C_{ws} = 0$

being trees, Cayley trees do not have triangles

## 7.6 POISSON NETWORKS

$$P(N) = \frac{e^c}{N} \quad c = \langle k \rangle$$

In Poisson networks the # of cycles of any size is finite when  $N \rightarrow \infty$ .

So such networks are "locally tree-like"

↓ This is the reason why we expect similar results to those obtained for Cayley trees

### PROPOSITION

The average distance  $\bar{l} = \langle d_{ij} \rangle$  of a Poisson network with average degree  $\langle k \rangle = e$  is

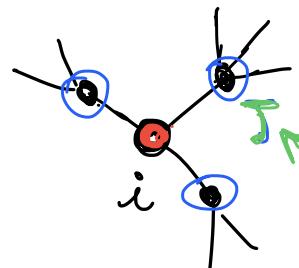
$$\boxed{\bar{l} \approx \frac{\ln N}{\ln e} \quad \text{for } N \gg 1}$$

Proof

(a)

# of nodes at distance  $d$  from a node  $i$

$$N_d(i) \approx \begin{cases} 1 & d=0 \\ k_i e^{d-1} & d>0 \end{cases}$$



- In 1 step from  $i$  we can reach  $k_i$  nodes

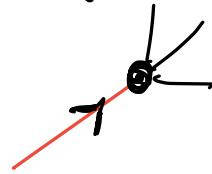
- Then the # of nodes that we can reach

from this branch (from node  $i$ ) is

$$b_j = k_{j-1}$$

branching ratio of node  $i$

average branching  
ratio of a node reached  
by following a link



$$\bar{b} = \frac{\langle k(k-1) \rangle}{\langle k \rangle} = \frac{c^2}{c} = c$$

See W 12

for POISSON  
NETWORK

$$N_d(i) \approx k_i c^{d-1}$$

Hence the average # of nodes at distance  $d$  from a RANDOMLY CHOSEN

node is

$$N_d \approx \begin{cases} 1 & d=0 \\ \langle k \rangle c^{d-1} = c^d & d>0 \end{cases}$$

(b)

$$N \simeq N_{d=l}$$

↑  
characteristic  
path length

$$N = e^l$$

$$\ln N = l \ln c$$

$$l \simeq \frac{\ln N}{\ln c} = \frac{\ln N}{\ln \langle k \rangle}$$

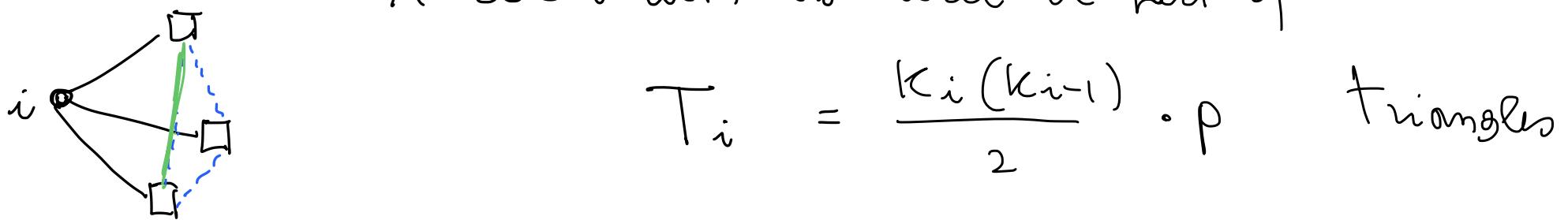
## PROPOSITION

The clustering coefficient  $C_{ws}$  of a Poisson network is

$$C_{ws} \approx \frac{c}{N} = \frac{\langle k \rangle}{N} = p \quad \text{for } N \gg 1$$

Proof

A node  $i$  with  $k_i$  will be part of



$$C_i = \frac{T_i}{\frac{k_i(k_i-1)}{2}} = p$$

$$C = p = \frac{c}{N}$$