## MTH4104 Cheat Sheet

## Shu Sasaki

## 13th February 2024

## 1 Chapter 2 and Chapter 3 (Week 1-3)

**GOAL**: Get used to an axiomatic approach to mathematics– given definitions/axioms, derive general statements about integers (that we know too well) via proofs and careful inspection of definitions etc.

**Proposition 1**. Let *a* and *b* be integers and suppose b > 0. Then a = bq + r for some integers *q* and  $0 \le r < b$ . The pair (q, r) is unique.

**Definition**. Let *a* and *b* be integers. We say that *a* divides *b* if there exists an integer *c* such that b = ac.

**Remark**. The only integer 0 divides is 0 itself.

**Definition**. Let *a* and *b* be integers. A common divisor of *a* and *b* is a non-negative integer *s* such that *s* divides both *a* and *b*. A gcd of *a* and *b* is the common divisor *r* satisfying the property that if *s* is another (different) common divisor of *a* and *b*, then s < r.

**Proposition 2**. *s* divides *r*.

We can say something similar for the lcm of *a* and *b*.

**Proposition 4**. If *a* is a non-negative integer, gcd(a, 0) = a. This is not a definition.

Lemma 5. gcd(a, b) = gcd(-a, b) = gcd(a, -b) = gcd(-a, -b). This is not a definition.

**Theorem 7 (Bezout's identity)**. Let *a* and *b* be integers. Then there exist integers *r* and *s* such that ar + bs = gcd(a, b).

The proof of Bezout explains only that these integers r and s exist and does not shed any light on how to actually find them. In practice, we make appeal to Euclid's algorithm instead.

Euclid's algorithm is based on the following proposition:

**Proposition 6.** Let *a* and *b* be integers. Suppose b > 0. By Proposition 1, there exists a unique pair of integers *q* and  $0 \le r < b$  such that a = bq + r. Then gcd(a, b) = gcd(b, r).

How do we use Euclid's algorithm to find *r* and *s* satisfying ar + bs = gcd(a, b)?

(NON-EXAMINABLE) If your Euclid's algorithm looks like:

$$\begin{array}{rcrcrcrc} \vdots \\ (s_n) & r_{n-2} & = & r_{n-1}q_n + r_n \\ (s_{n+1}) & r_{n-1} & = & r_nq_{n+1} + r_{n+1} \\ & \vdots \\ (s_N) & r_{N-1} & = & r_Nq_{N+1} + r_{N+1} \\ (s_{N+1}) & r_N & = & r_{N+1}q_{N+2} \end{array}$$

then we know that gcd(a, b) is  $r_{N+1}$ , because we may repeat Proposition 6 to deduce that

$$\gcd(a,b) = \cdots = \gcd(r_{n-2},r_{n-1}) \stackrel{(s_n)}{=} \gcd(r_{n-1},r_n) \stackrel{(s_{n+1})}{=} \gcd(r_n,r_{n+1}) = \cdots = \\ \gcd(r_{N-1},r_N) \stackrel{(s_N)}{=} \gcd(r_N,r_{N+1}) \stackrel{(s_{N+1})}{=} r_{N+1}.$$

We also see from  $(s_N)$  that  $r_{N+1} = -q_{N+1}r_N + r_{N-1}$ . Indeed, for every n (e.g.  $N, N-1, \ldots$ ), there exist integers  $X_n$  and  $Y_n$  satisfying

$$r_{N+1} = X_n r_n + Y_n r_{n-1}.$$

This will find us *r* and *s* such that  $ar + bs = r_{N+1}$ .

We may prove the assertion by induction 'in reverse' (one can reindex all to make this rigorous). We saw  $(X_N, Y_N) = (-q_N, 1)$  does the job. Supposing that there exist integers  $X_n$  and  $Y_n$  such that

$$r_{N+1} = X_n r_n + Y_n r_{n-1}$$

we aim at proving that there exists  $X_{n-1}$  and  $Y_{n-1}$  such that

$$r_{N+1} = X_{n-1}r_{n-1} + Y_{n-1}r_{n-2}.$$

We will spell out  $X_{n-1}$  and  $Y_{n-1}$  in terms of  $X_n$  and  $Y_n$ . To see this, plug  $r_n = (-q_n)r_{n-1} + r_{n-2}$  obtained from  $(s_n)$  into  $r_{N+1} = X_n r_n + Y_n r_{n-1}$ . We then get

$$r_{N+1} = X_n((-q_n)r_{n-1} + r_{n-2}) + Y_nr_{n-1} = (-q_nX_n + Y_n)r_{n-1} + X_nr_{n-2},$$

hence  $(X_{n-1}, Y_{n-1}) = (-q_n X_n + Y_n, X_n)$  does the job. It is possible to use this inductively (as *n* decreases) to find *X*'s and *Y*'s, starting with  $(X_N, Y_N) = (-q_N, 1)$ .

**Definition**. A prime number is a positive integer *n* whose positive integer divisor is 1 or itself. Alternatively, we may define it as a positive integer whose integer divisors are  $\{\pm 1, \pm n\}$ .

By Bezout, this is equivalent to the following: if a and b are integers and n divides ab, then n divides either a or b. The latter definition allows us to prove:

Theorem 8 (the Fundamental Theorem of Arithmetic). Every integer is of the form

$$(-1)^{r_{\infty}}\prod_{p}p^{r_{p}}$$

for some non-negative integers  $r_{\infty}$  and  $r_p$ , up to reordering of prime factors. The power  $r_p$  is the maximum number of time p divides the integer. For example,  $45 = 3^2 \cdot 5$  so  $r_p = 0$  if p is not 3 nor 5,  $r_3 = 2$ ,  $r_5 = 1$  and  $r_{\infty} = 0$ .

Let  $\mathcal{R}$  be a relation on S. We let  $[a] = [a]_{\mathcal{R}}$  denote the subset of all b in S which are related to a, i.e.  $a\mathcal{R}b$ . If  $\mathcal{R}$  is an equivalence relation (satisfying a set of conditions), then

$$a\mathcal{R}b$$
 if and only if  $[a] = [b]$ .

**Theorem 9**. Given a set *S*, there exists a bijective correspondence between

- the equivalence relations  $\mathcal{R}$  on S,
- the partitions  $\mathscr{P}$  (a set of subsets of S satisfying certain conditions) on S.

**Proposition 10**. Let *n* be a positive integer. Then  $(\mathcal{R}, S) = (\equiv \mathbb{Z})$ , defined such that  $a \equiv b \mod n$  if and only if *n* divides b - a (for integers *a* and *b*), is an equivalence relation.

**Definition**. Let  $\mathbb{Z}_n$  denote the set of equivalence classes [a] with respect to  $(\equiv, \mathbb{Z})$ .

Since  $a \equiv b \mod n$  if and only if [a] = [b], a lot of equivalence classes may be identified. Indeed,

**Proposition 11**.  $|\mathbb{Z}_n| = n$ .

Proposition 1 proves Proposition 11. Indeed, if *a* is an integer (*n* is, by definition, a positive integer), then there exists *q* and  $0 \le r < n$  such that a = nq + r. Therefore  $a \equiv r$ , i.e. [a] = [r]. The proof also elaborates that  $\mathbb{Z}_n = \{[0], [1], \ldots, [n-1]\}$ . The element [r] is nothing other than the set of integers *b* with remainder *r* when divided by *n* (i.e.  $b \equiv r \mod n$ ).

On  $\mathbb{Z}_n$ , we define  $+, -, \times$ :

$$\begin{array}{rcl} [a] + [b] &=& [a+b] \\ [a] - [b] &=& [a-b] \\ [a] [b] &=& [ab] \end{array}$$

but no division. These do not depend on choice of representatives, i.e. if  $a \equiv a' \mod n$ , then [a] + [b] = [a'] + [b] etc.

No division is defined but:

**Definition**. We say that [a] of  $\mathbb{Z}_n$  has multiplicative inverse if there exists an integer b such that [a][b] = [1] (or equivalently  $ab \equiv 1 \mod n$ ). This plays the role of 1/[a] but not literally (1/[a] or [1/a] simply does not make sense!). The multiplicative inverse is often written as  $[a]^{-1}$ .

**Remark**. The multiplicative inverse, if exists, is unique. Suppose that [b] and [c] are elements of  $\mathbb{Z}_n$  such that [a][b] = [1] and [a][c] = [1]. Multiplying both sides of [c][a] = [1] by [b], we obtain [c][a][b] = [1][b], i.e. [c] = [b].

**Theorem 12**. An element [a] of  $\mathbb{Z}_n$  has multiplicative inverse if and only if gcd(a, n) = 1.

The proof explains how to find the multiplicative inverse explicitly. If a is an integer such that gcd(a, n) = 1 (which one can check in practice by Euclid's algorithm), Euclid's algorithm find integers b and c such that ab + nc = gcd(a, n) = 1. It then follows that  $ab \equiv 1 \mod n$ , i.e. [a][b] = [ab] = [1].

**Proposition 13**. An element [a] of  $\mathbb{Z}_n$  has no multiplicative inverse if and only if there exists b, not congruent to 0 mod n, such that [a][b] = [0].

Example.  $[2]_6[3]_6 = [0]_6$ .

It is possible to compute the number of elements in  $\mathbb{Z}_n$  with multiplicative inverses, using the fundamental theorem of arithmetic: if  $=\prod_{p} p^{r_p}$ , then it is computed by  $\prod_p (p-1)p^{r_p-1}$ .

What is it useful for? It is possible to solve 'linear congruence equations':  $ax + b \equiv c \mod n$ (when gcd(a, n) = 1). Indeed,  $[x] = [c - b][a]^{-1}$  where  $[a]^{-1}$  is the multiplicative inverse of [a](this is NOT 1/[a]). What if gcd(a, n) > 1? Take Number Theory next year!