

Linear algebra

Throughout work in \mathbb{R}^n (n -dimensional Euclidean space)

If $\underline{x} \in \mathbb{R}^n$ then $\underline{x} = \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix}$ where $x_1, x_2, \dots, x_n \in \mathbb{R}$ *entries*

e.g. $\underline{y} = \begin{pmatrix} 2 \\ 4 \\ 3 \end{pmatrix} \in \mathbb{R}^3$ where $y_1=2, y_2=4, y_3=3$

$\underline{0} = \begin{pmatrix} 0 \\ 0 \\ \vdots \\ 0 \end{pmatrix}$ zero vector (dimension clear from context)

If A $n \times m$ matrix (n rows, m columns)

e.g. $A = \begin{pmatrix} 2 & 3 & 1 \\ 1 & 4 & 5 \end{pmatrix}$ Then $A^T = \begin{pmatrix} 2 & 1 \\ 3 & 4 \\ 1 & 5 \end{pmatrix}$

Convention

$\underline{x} = \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix} \in \mathbb{R}^n$ is a $n \times 1$ matrix (column vector)

$\underline{x}^T = (x_1, x_2, \dots, x_n)$ $1 \times n$ matrix (row vector)

$\underline{z} = \begin{pmatrix} z_1 \\ z_2 \\ \vdots \\ z_n \end{pmatrix} \in \mathbb{R}^n$

$\underline{x}^T \underline{z} = x_1 z_1 + x_2 z_2 + \dots + x_n z_n$
 $= \sum x_i z_i$ (called dot product or scalar product)

For $\underline{x}_1, \underline{x}_2, \dots, \underline{x}_n \in \mathbb{R}^m$ and $a_1, a_2, \dots, a_n \in \mathbb{R}$

$$a_1 \underline{x}_1 + a_2 \underline{x}_2 + \dots + a_n \underline{x}_n \in \mathbb{R}^m$$

is called linear combination of vectors $\underline{x}_1, \dots, \underline{x}_n$

For $m=1$, $x_1, \dots, x_n \in \mathbb{R}$ are real numbers or

linear combination for real numbers or variables
is given by

$$a_1 x_1 + a_2 x_2 + \dots + a_n x_n$$

can be written concisely as

$$\underline{a}^T \underline{x} \quad \text{where} \quad \underline{a} = \begin{pmatrix} a_1 \\ \vdots \\ a_n \end{pmatrix} \quad \underline{x} = \begin{pmatrix} x_1 \\ \vdots \\ x_n \end{pmatrix}.$$

For constants $a_1, a_2, \dots, a_n \in \mathbb{R}$, $b \in \mathbb{R}$

- (1) $a_1 x_1 + a_2 x_2 + \dots + a_n x_n = b$ ← linear equation over variables x_1, \dots, x_n
- (2) $a_1 x_1 + a_2 x_2 + \dots + a_n x_n \leq b$ ← linear inequality over variables x_1, \dots, x_n

Can write this concisely as

(1) $\underline{a}^T \underline{x} = b$

(2) $\underline{a}^T \underline{x} \leq b$

For $\underline{x} = \begin{pmatrix} x_1 \\ \vdots \\ x_n \end{pmatrix}$ and $\underline{y} = \begin{pmatrix} y_1 \\ \vdots \\ y_n \end{pmatrix} \in \mathbb{R}^n$

Write $\underline{x} \geq \underline{y}$ to mean $x_1 \geq y_1, x_2 \geq y_2, \dots, x_n \geq y_n$

For $\underline{x} = \begin{pmatrix} x_1 \\ \vdots \\ x_n \end{pmatrix}$ and $\underline{y} = \begin{pmatrix} y_1 \\ \vdots \\ y_n \end{pmatrix} \in \mathbb{R}^n$

Write $\underline{x} \succeq \underline{y}$ to mean $x_1 \succeq y_1, x_2 \succeq y_2, \dots, x_n \succeq y_n$

True or false $\begin{pmatrix} 8 \\ 2 \\ 4 \end{pmatrix} \succeq \begin{pmatrix} 7 \\ 2 \\ 3 \end{pmatrix}$? True

If $\underline{x}, \underline{y} \in \mathbb{R}^n$ then $\underline{x} \succeq \underline{y}$ or $\underline{y} \succeq \underline{x}$ False e.g. $\begin{pmatrix} 2 \\ 1 \end{pmatrix} \begin{pmatrix} 1 \\ 2 \end{pmatrix}$

If $\underline{x}, \underline{y}, \underline{z} \in \mathbb{R}^n$ and $\underline{x} \succeq \underline{y}$ and $\underline{y} \succeq \underline{z}$ then $\underline{x} \succeq \underline{z}$ True.

Suppose A is an $m \times n$ matrix

$$A = \begin{pmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ a_{21} & a_{22} & \dots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ a_{m1} & a_{m2} & \dots & a_{mn} \end{pmatrix} \quad \underline{x} = \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix} \in \mathbb{R}^n$$

$$\underline{b} = \begin{pmatrix} b_1 \\ b_2 \\ \vdots \\ b_m \end{pmatrix} \in \mathbb{R}^m$$

$$A\underline{x} = \underline{b}$$

Suppose A is an $m \times n$ matrix

$$A = \begin{pmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ a_{21} & a_{22} & \dots & a_{2n} \\ \vdots & & & \\ a_{m1} & a_{m2} & \dots & a_{mn} \end{pmatrix} \quad \underline{x} = \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix} \in \mathbb{R}^n$$

$$\underline{b} = \begin{pmatrix} b_1 \\ b_2 \\ \vdots \\ b_m \end{pmatrix} \in \mathbb{R}^m$$

$A\underline{x} = \underline{b}$ is a concise way of writing

$$a_{11}x_1 + a_{12}x_2 + \dots + a_{1n}x_n = b_1$$

$$a_{21}x_1 + a_{22}x_2 + \dots + a_{2n}x_n = b_2$$

\vdots

$$a_{m1}x_1 + a_{m2}x_2 + \dots + a_{mn}x_n = b_m$$

$A\underline{x} \leq \underline{b}$ is a concise way of writing

$$a_{11}x_1 + a_{12}x_2 + \dots + a_{1n}x_n \leq b_1$$

$$a_{21}x_1 + a_{22}x_2 + \dots + a_{2n}x_n \leq b_2$$

\vdots

$$a_{m1}x_1 + a_{m2}x_2 + \dots + a_{mn}x_n \leq b_m$$

Example of mathematical program

↑ **goal**

maximise $x_1^2 - \sqrt{x_2} + \log(x_1 + x_3)$ ← **objective function**

subject to

$$\left. \begin{aligned} x_1 + x_3^2 &\leq 2 \\ x_2 + x_1 x_3 &\geq 5 \\ x_1 + x_3 &= 3 \end{aligned} \right\} \text{constraints}$$

$x_1 \geq 0, x_2 \leq 0, x_3$ unrestricted

Mark the different elements

↑ **variables with sign restrictions**

Defn

A mathematical program is an optimisation problem consisting of 4 components:

- 1) A set of variables with possible sign restrictions (always listed at the end)
- 2) An objective function (i.e. a function of the variables that we want to optimise)
- 3) A goal (i.e. to maximise or minimise objective function)
- 4) A set of constraints on the variables given by equations or inequalities.

Defn A linear program is a mathematical program where

- variables assumed to be continuous and real
(e.g. cannot constrain variable to be integer)
- the objective function is a linear combination of the variables
- each constraint is a linear inequality or linear equation over the variables.

Examples of linear program

Ex1 maximise $2x_1 - 5x_2 + 3x_3$
subject to $x_1 + x_2 + x_3 \leq 7$
 $x_1 + 3x_3 \leq 5$
 $x_1, x_2, x_3 \geq 0$

Ex2 minimise $3x_1 - x_2$
subject to $-x_1 + 6x_2 + x_3 \geq -3$
 $x_1 + x_2 = 1$
 $x_3 \leq 2$

x_1 unrestricted
 $x_2 \geq 0$
 $x_3 \leq 0$

Quiz: which of these are linear programmes and which are not?

Which of these is a linear program?

Maximise $2x + 3y - 10z$
Subject to $x \leq 3$

$$4 \leq 2y - z \leq 5 + 3x$$

$x \geq 0, y \leq 0,$
 z unrestricted



can write as two separate inequalities, namely

$$4 \leq 2y - z$$
$$2y - z \leq 5 + 3x$$

Then rearrange these:

$$2y - z \geq 4$$
$$-3x + 2y - z \leq 7$$

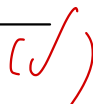
replace ↙

Minimise $(2x_1 + 1)^2 + 2x_2 - 4x_1^2 - 1$

subject to $x_1 + x_2 \geq 3$

$$x_1 - x_2 \leq 5$$

$$x_1, x_2 \geq 0$$



After expanding objective function becomes

$$4x_1 + 2x_2$$

maximise $2x_1 + 3x_2 - x_3 + 6$

subject to $x_1 + 2x_2 \geq 3$

$$x_1 + 5x_2 = 10$$

x_1, x_2, x_3 unrestricted



objective function not linear combination because of 6.

Defn We say that a linear program is in standard (inequality) form if it is written as

$$\begin{aligned} & \text{maximise} && \underline{c}^T \underline{x} \\ & \text{Subject to} && A \underline{x} \leq \underline{b} \\ & && \underline{x} \geq \underline{0} \end{aligned}$$

Where A is an $m \times n$ matrix of real numbers

$$\underline{c} \in \mathbb{R}^n, \quad \underline{b} \in \mathbb{R}^m, \quad \underline{x} = \begin{pmatrix} x_1 \\ \vdots \\ x_n \end{pmatrix} \text{ is vector of variables.}$$

Expand above: assume that

$$A = \begin{pmatrix} a_{11} & \dots & a_{1n} \\ \vdots & \ddots & \vdots \\ a_{m1} & \dots & a_{mn} \end{pmatrix} \quad \underline{c} = \begin{pmatrix} c_1 \\ \vdots \\ c_n \end{pmatrix} \quad \underline{b} = \begin{pmatrix} b_1 \\ \vdots \\ b_m \end{pmatrix}$$

Above becomes

$$\begin{aligned} & \text{maximise} && c_1 x_1 + c_2 x_2 + \dots + c_n x_n \\ & \text{Subject to} && a_{11} x_1 + a_{12} x_2 + \dots + a_{1n} x_n \leq b_1 \\ & && a_{21} x_1 + a_{22} x_2 + \dots + a_{2n} x_n \leq b_2 \\ & && \vdots \\ & && a_{m1} x_1 + a_{m2} x_2 + \dots + a_{mn} x_n \leq b_m \\ & && x_1, x_2, \dots, x_n \geq 0 \end{aligned}$$

Why do we want to do this?

Write Ex1 and Ex2 in standard inequality form

$$\begin{array}{ll} \text{maximise} & \underline{c}^T \underline{x} \\ \text{Subject to} & A \underline{x} \leq \underline{b} \\ & \underline{x} \geq \underline{0} \end{array}$$

$$\begin{array}{ll} \underline{\text{Ex1}} \text{ maximise} & 2x_1 - 5x_2 + 3x_3 \\ \text{Subject to} & x_1 + x_2 + x_3 \leq 7 \\ & x_1 + 3x_3 \leq 5 \\ & x_1, x_2, x_3 \geq 0 \end{array}$$

$$\begin{array}{ll} \text{maximise} & (2, -5, 3) \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} \\ \text{subject to} & \begin{pmatrix} 1 & 1 & 1 \\ 1 & 0 & 3 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} \leq \begin{pmatrix} 7 \\ 5 \end{pmatrix} \\ & \underline{x} \geq \underline{0} \end{array}$$

A \underline{x} \underline{b}

What about ex2. Have some problems

$$\begin{array}{ll} \text{maximise} & \underline{c}^T \underline{x} \\ \text{Subject to} & A \underline{x} \leq \underline{b} \\ & \underline{x} \geq \underline{0} \end{array}$$

$$\begin{array}{ll} \underline{\text{Ex2}} \text{ minimise} & 3x_1 - x_2 \\ \text{Subject to} & -x_1 + 6x_2 + x_3 \geq -3 \\ & x_1 + x_2 = 1 \\ & x_3 \leq 2 \end{array}$$

$$\begin{array}{l} x_1 \text{ unrestricted} \\ x_2 \geq 0 \\ x_3 \leq 0 \end{array}$$

How to transform any linear program to standard inequality form

① For each variable x_i if sign constraint is

$$x_i \geq 0 \quad \checkmark$$

$$x_i \leq 0 \quad \text{replace } x_i \text{ with } \bar{x}_i \geq 0 \quad \text{where } x_i = -\bar{x}_i$$

$$x_i \text{ unrestricted} \quad \text{replace } x_i \text{ with } x_i^+ - x_i^-$$

with $x_i^+ \geq 0 \quad x_i^- \geq 0$

② If goal is $\min \underline{c}^T \underline{x}$ replace with $\max (-\underline{c})^T \underline{x}$

minimizing a fn is same as maximising negative of that function

③ For each constraint, if constraint is

$$\underline{a}^T \underline{x} \leq b \quad \checkmark$$

$$\underline{a}^T \underline{x} \geq b \quad \text{replace with } (-\underline{a})^T \underline{x} \leq -b$$

$$\underline{a}^T \underline{x} = b \quad \text{replace with two constraints}$$

$$\underline{a}^T \underline{x} \leq b$$

$$(-\underline{a})^T \underline{x} \leq -b$$

Apply to Ex2

e.g. $r \leq 2 \Leftrightarrow r = 2$
 $-r \leq -2$

①

Ex2 minimise $3x_1 - x_2$
 subject to $-x_1 + 6x_2 + x_3 \geq -3$
 $x_1 + x_2 = 1$
 $x_3 \leq 2$

x_1 unrestricted
 $x_2 \geq 0$
 $x_3 \leq 0$

min $3x_1^+ - 3x_1^- - x_2$
 subjc $-x_1^+ + x_1^- + 6x_2 - \bar{x}_3 \geq -3$
 $x_1^+ - x_1^- + x_2 = 1$
 $-\bar{x}_3 \leq 2$

$$x_1^+, x_1^-, x_2, \bar{x}_3 \geq 0$$

$$x_1 = x_1^+ - x_1^-$$

$$x_3 = -\bar{x}_3$$

① Fixed sign restrictions

$$\begin{aligned} \min \quad & 3x_1^+ - 3x_1^- - x_2 \\ \text{subtc} \quad & -x_1^+ + x_1^- + 6x_2 - \bar{x}_3 \geq -3 \\ & x_1^+ - x_1^- + x_2 = 1 \\ & -\bar{x}_3 \leq 2 \end{aligned}$$

$$x_1^+, x_1^-, x_2, \bar{x}_3 \geq 0$$

$$\begin{aligned} x_1 &= x_1^+ - x_1^- \\ x_3 &= -\bar{x}_3 \end{aligned}$$

② Fix goal

$$\begin{aligned} \max \quad & -3x_1^+ + 3x_1^- + x_2 \\ \text{subtc} \quad & -x_1^+ + x_1^- + 6x_2 - \bar{x}_3 \geq -3 \\ & x_1^+ - x_1^- + x_2 = 1 \\ & -\bar{x}_3 \leq 2 \end{aligned}$$

$$x_1^+, x_1^-, x_2, \bar{x}_3 \geq 0$$

$$\begin{aligned} x_1 &= x_1^+ - x_1^- \\ x_3 &= -\bar{x}_3 \end{aligned}$$

③ Fix constraint

$$\begin{aligned} \max \quad & -3x_1^+ + 3x_1^- + x_2 \\ \text{subtc} \quad & +x_1^+ - x_1^- - 6x_2 + \bar{x}_3 \leq +3 \\ & x_1^+ - x_1^- + x_2 \leq 1 \\ & -x_1^+ + x_1^- - x_2 \leq -1 \\ & -\bar{x}_3 \leq 2 \end{aligned}$$

$$x_1^+, x_1^-, x_2, \bar{x}_3 \geq 0$$

$$\begin{aligned} x_1 &= x_1^+ - x_1^- \\ x_3 &= -\bar{x}_3 \end{aligned}$$



Given optimal solution for transformed program, can use this substitution to get optimal solution to original.

In matrix form $\max c^T x$

$$\text{subtc } Ax \leq b, \quad x \geq 0$$

$$x = \begin{pmatrix} x_1^+ \\ x_1^- \\ x_2 \\ \bar{x}_3 \end{pmatrix}$$

$$c = \begin{pmatrix} -3 \\ 3 \\ 1 \\ 0 \end{pmatrix}$$

$$A = \begin{pmatrix} 1 & -1 & -6 & 1 \\ 1 & -1 & 1 & 0 \\ -1 & 1 & -1 & 0 \\ 0 & 0 & 0 & -1 \end{pmatrix}$$

$$b = \begin{pmatrix} 3 \\ 1 \\ -1 \\ 2 \end{pmatrix}$$

Careful about order of variables in vector. If we change order then A and c change.

Remark As a convention, when ordering variables,

- put variables in alphabetical order
- lower indices before higher indices
- '+'s before '-'s

e.g. $x_1, x_2^+, x_2^-, x_3, x_4, y_1^+, y_1^-, y_2$

The next two pages cover further revision of linear algebra that will be used later in the module.

This was not covered in week 1 but it will help you to do the seminar questions

Matrix multiplication interpreted as linear combinations

When $\underline{x}_1, \dots, \underline{x}_n \in \mathbb{R}^m$ and $a_1, \dots, a_n \in \mathbb{R}$ we have

$$a_1 \underline{x}_1 + a_2 \underline{x}_2 + \dots + a_n \underline{x}_n = \begin{pmatrix} \underline{x}_1 & \underline{x}_2 & \dots & \underline{x}_n \end{pmatrix} \begin{pmatrix} a_1 \\ a_2 \\ \vdots \\ a_n \end{pmatrix}$$

$m \times n$ matrix whose
n columns are the
vectors $\underline{x}_1, \dots, \underline{x}_n$ $\quad = \quad \underline{X} \underline{a}$

Linear independence and rank. (will be used in later weeks)

We say vectors $\underline{a}_1, \dots, \underline{a}_m \in \mathbb{R}^n$ are linearly dependent if $\exists x_1, \dots, x_m$ not all zero s.t.

$$x_1 \underline{a}_1 + x_2 \underline{a}_2 + \dots + x_m \underline{a}_m = \underline{0}$$

i.e. $\begin{pmatrix} \underline{a}_1 & \underline{a}_2 & \dots & \underline{a}_m \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_m \end{pmatrix} = \underline{0}$

call this
 $n \times m$ matrix A

call this vector $\underline{x} \in \mathbb{R}^m$

i.e. $A \underline{x} = \underline{0}$ for some $\underline{x} \neq \underline{0}$

Show $\begin{pmatrix} 1 \\ 1 \\ 0 \end{pmatrix}, \begin{pmatrix} 1 \\ 0 \\ 1 \end{pmatrix}, \begin{pmatrix} 2 \\ 1 \\ 1 \end{pmatrix}$ are linearly dependent.

Say $\underline{a}_1, \dots, \underline{a}_m$ are linearly independent if they are not linearly dependent.

So $\underline{a}_1, \dots, \underline{a}_m$ are linearly independent if,

whenever $x_1 \underline{a}_1 + x_2 \underline{a}_2 + \dots + x_m \underline{a}_m = \underline{0}$,

we have $x_1 = x_2 = \dots = x_m = 0$

i.e. if $A \underline{x} = \underline{0}$ has unique solution $\underline{x} = \underline{0}$

Rank of a matrix $m \times n$

For $m \times n$ matrix A , its n columns are vectors in \mathbb{R}^m
its m rows are vectors in \mathbb{R}^n

column rank(A) = $\max \{ k : A \text{ has } k \text{ linearly independent columns} \}$

row rank(A) = $\max \{ l : A \text{ has } l \text{ linearly independent rows} \}$

Fact row rank(A) = column rank(A)

Just call it rank(A).

$$\text{rank}(A) \leq \min(m, n)$$

Fact If A is $n \times n$ matrix then

$$\text{rank}(A) = n \iff A \text{ is invertible}$$

$$\iff A\underline{x} = \underline{b} \text{ has unique solution for every } \underline{b} \in \mathbb{R}^n.$$